





0x0030:	0000	0000	0000	0000	0000	0000	0000	0000	.....
0x0040:	0000	0000	0000	0000	0000	0000	0000	0000	.....
0x0050:	0000	0000	0000	0000	0000	0000	0000	0000	.....
0x0060:	0000	0000	0000	0000	0000	0000	0000	0000	.....
0x0070:	0000	0000	0000	0000	0000	0000	0000	0000	.....
0x0080:	0000	0000	0000	0000	0000	0000	515c	0000	.....Q\..
0x0090:	b8f5	0909	d843	2309	0000	0000	0000	0000	....C#.....
0x00a0:	b81a	a040	3500	0d08	2000	0000	2000	0000	...@5.....
0x00b0:	0000	0000	0100	0000	00e3	2309	1704	0000	.....#.....
0x00c0:	b81a	a141	3500	0000	20d0	2309	8100	0000	...A5.....#.....
0x00d0:	2811	2409	48cf	2309	20e7	2309	1704	0000	(.\$.H.#...#.....
0x00e0:	c16c	5b9a	3500	0035	40d0	2309	6100	0000	.l[.5..5@.#.a... Hj"...#.@.#.....
0x00f0:	486a	2209	98d6	2309	40eb	2309	1704	0000	...C5.....
0x0100:	173d	c743	3500	0000	2000	0000	2000	0000	.....`#.....
0x0110:	0000	0000	0100	0000	60ef	2309	1704	0000	...d.B5.....#!... _Q4.h(\$...#.....
0x0120:	5f64	ae42	3500	0000	80d0	2309	2100	0000	`..1C5.....
0x0130:	a851	3409	6828	2409	80f3	2309	1704	0000	.....
0x0140:	6007	3143	3500	0000	c000	0000	9000	0000	.....
0x0150:	0f00	0000	0000	0000	0000	0000	0000	0000	.....
0x0160:	501b	f2f4	0100	0000	0000	0000	0000	0000	P.....
0x0170:	0000	0000	0000	0000	0000	0000	0000	0000	.....
0x0180:	0000	0000	0000	0000	0000	0000	0000	0000	.....
0x0190:	0000	0000	0000	0000	0000	0000	0000	0000	.....
0x01a0:	0000	0000	0000	0000	0000	0000	0000	0000	.....
0x01b0:	0000	0000	0000	0000	0000	0000	0000	0000	.....
0x01c0:	0000	0000	0000	0000	0000	0000	0000	0000	.....
0x01d0:	0000	0000	0000	0000	0000	0000	8100	0000	.....
0x01e0:	f04a	0d08	6828	2409	a0f7	2309	1704	0000	.J..h(\$...#..... ...@5.h.....
0x01f0:	b81a	a040	3500	6809	2000	0000	2000	0000	...#.....#.....
0x0200:	a8d1	2309	0100	0000	c0fb	2309	1704	0000	...A5...@.....
0x0210:	b81a	a141	3500	0000	4000	0000	2000	0000	...#.....#.....
0x0220:	c8d1	2309	0100	0000	e0ff	2309	1704	0000	.l[.5...`..... ..\$....\$......
0x0230:	c16c	5b9a	3500	0000	6000	0000	2000	0000	...C5.....
0x0240:	1808	2409	0100	0000	0004	2409	1704	0000	.....
0x0250:	173d	c743	3500	0000	d001	0000	9000	0000	...\$.....
0x0260:	0f00	0000	0000	0000	8019	0000	9000	0000	.....
0x0270:	0f00	0000	0000	0000	0000	0000	0000	0000	.....
0x0280:	50fb	5fed	0100	0000	0000	0000	0000	0000	P_.....
0x0290:	0000	0000	0000	0000	0000	0000	0000	0000	.....
0x02a0:	0000	0000	0000	0000	0000	0000	0000	0000	.....
0x02b0:	0000	0000	0000	0000	0000	0000	0000	0000	.....
0x02c0:	0000	0000	0000	0000	0000	0000	0000	0000	.....
0x02d0:	0000	0000	0000	0000	0000	0000	0000	0000	.....
0x02e0:	0000	0000	0000	0000	0000	0000	0000	0000	.....
0x02f0:	0000	0000	0000	0000	0000	0000	8119	0000	.....
0x0300:	f8f2	2209	b820	2409	0000	0000	0000	0000	...".\$......
0x0310:	b81a	a040	3500	2309	0000	0000	6119	0000	...@5.#.....a... x... (w#.....
0x0320:	78f6	1909	2877	2309	0000	0000	0000	0000	...@...5...vD... X...u.....
0x0330:	b81a	a040	fd00	0035	fd00	7644	0000	0000	X....~.....
0x0340:	5802	ffff	b075	0000	0204	05b4	0101	0402	X.... .....
0x0350:	5802	ffff	cc7e	0000	0204	05b4	0101	0402	P...'......
0x0360:	5802	ffff	b87c	0000	0204	05b4	0101	0402	.D#...#.....A... H..."RL.....
0x0370:	5002	ffff	b327	0000	0204	05b4	0101	0402	_d.B28.528d8... X...\$......
0x0380:	b844	2309	f8b8	2309	0000	0000	4108	0000	...C5.....#!... ..#.(ST.....
0x0390:	48e0	2209	c852	4c09	0000	0000	0000	0000	...C...5..SP... P.....
0x03a0:	5f64	ae42	3238	0035	3238	6438	0000	0000	.....
0x03b0:	58c2	ffff	ef24	0000	0204	05b4	0101	0402	.....
0x03c0:	173d	c743	3500	0000	00cf	2309	2104	0000	.....
0x03d0:	98d7	2309	2853	5409	0000	0000	0000	0000	.....
0x03e0:	173d	c743	0420	0035	0420	5350	0000	0000	.....
0x03f0:	5002	ffff	e180	0000	0204	05b4	0101	0402	.....
0x0400:	0000	0000	0000	0000	0000	0000	0000	0000	.....
0x0410:	0000	0000	0000	00					.....

Figure 3: Captured sample #2

We see a 100% duplicate payload, confirming the spoofing capabilities explained in section 4 of the Xor advisory.

The two packets are generated by the same machine while spoofing only the last octet of its real public IP address. The payload consists of garbage memory data of what happens to be in the used memory space. The memory can also contain passwords and ssh private keys.

In the DNS Flood attack payloads, we see all distinguishable characteristics as well.

```
08:03:32.375073 IP xx.xx.223.184.21324 > 96.7.49.67.53: 45996 [1au] A? xx.xxxx20.com. (43)
  0x0000: 4500 0047 4c53 0000 4c11 64fd xxxx dfb8 E..GLS..L.d....
  0x0010: 6007 3143 534c 0035 0033 8fea b3ac 0020 \.1CSL.5.3.....
  0x0020: 0001 0000 0000 0001 02xx xx07 xxxx xxxx .....xx.xxxx
  0x0030: xx32 3003 636f 6d00 0001 0001 0000 2910 x20.com.....).
  0x0040: 0000 0000 0000 00
08:03:32.380600 IP xx.xx.223.184.21324 > 96.7.49.67.53: 45996 [1au] A? www.xxxxxx1.com. (44)
  0x0000: 4500 0048 4c53 0000 4c11 64fc xxxx dfb8 E..HLS..L.d....
  0x0010: 6007 3143 534c 0035 0034 33a3 b3ac 0020 \.1CSL.5.43.....
  0x0020: 0001 0000 0000 0001 0377 7777 07xx xxxx .....www.xxx
  0x0030: xxxx xx31 0363 6f6d 0000 0100 0100 0029 xxx1.com.....)
  0x0040: 1000 0000 0000 0000
08:03:32.392463 IP xx.xx.223.184.21324 > 96.7.49.67.53: 45996 [1au] A? www.xxxxxxxxxxxxxxx.com.
(51)
  0x0000: 4500 004f 4c53 0000 4c11 64f5 xxxx dfb8 E..OLS..L.d....
  0x0010: 6007 3143 534c 0035 003b e168 b3ac 0020 \.1CSL.5.;.h....
  0x0020: 0001 0000 0000 0001 0377 7777 0exx xxxx .....www.xxx
  0x0030: xxxx xxxx xxxx xxxx xxxx xx03 636f 6d00 xxxxxxxxxxxxx.com.
  0x0040: 0001 0001 0000 2910 0000 0000 0000 00
08:03:32.398063 IP xx.xx.223.184.21324 > 96.7.49.67.53: 45996 [1au] A? www.xxxxxx1.com. (44)
  0x0000: 4500 0048 4c53 0000 4c11 64fc xxxx dfb8 E..HLS..L.d....
  0x0010: 6007 3143 534c 0035 0034 33a3 b3ac 0020 \.1CSL.5.43.....
  0x0020: 0001 0000 0000 0001 0377 7777 07xx xxxx .....www.xxx
  0x0030: xxxx xx31 0363 6f6d 0000 0100 0100 0029 xxx1.com.....)
  0x0040: 1000 0000 0000 0000
08:03:32.415895 IP xx.xx.223.184.21324 > 96.7.49.67.53: 45996 [1au] A? xx.xxxxx20.com. (43)
  0x0000: 4500 0047 4c53 0000 4c11 64fd xxxx dfb8 E..GLS..L.d....
  0x0010: 6007 3143 534c 0035 0033 8ee0 b3ac 0020 \.1CSL.5.3.....
  0x0020: 0001 0000 0000 0001 02xx xx07 xxxx xxxx .....xx.xxxx
  0x0030: xx32 3003 636f 6d00 0001 0001 0000 2910 x20.com.....).
  0x0040: 0000 0000 0000 00
```

Figure 4: Sample DNS Flood

As you can see, the observed DNS query flood is sourcing from the same source IP, with the same source port and the same query ID targeting multiple domains. The structure of all the packets might also look similar. This is because all of the target domains are passed from the C&C to the bots in a single command. From there the infected machines build the same template for all of the target domains.

We also see the Additional Record Count bit set across all packets and the DNSSEC options at the end.

Another not-so-important characteristic we noticed is that sometimes the DNS queries include the RD (Recursion Desired) flag on and sometimes it's off.

**3.0 / RECOMMENDED DETECTION METHODS/** To detect this botnet in your network, one can look for the communications between bot and C2 using the following Snort rule:

```
alert TCP $HOME_NET any -> $EXTERNAL_NET any ( msg: "Xor-DDoS"; \
flow: established; \
content: "BB2FA36AAA9541F0BB2FA36AAA9541F0"; \
offset:32; depth: 64; \
classtype: trojan-activity; \
sid: 201500010; rev: 1;)
```

Figure 5: Snort rule matches on the XOR key string within the appropriate offsets in the initial registration between bot and C2

To detect infection of this malware on your hosts you can use the following Yara rule:

```
rule XorDDoSv1
{
  meta:
    author = "Akamai SIRT"
    description = "Rule to detect XorDDoS infection"

  strings:
    $st0 = "BB2FA36AAA9541F0"
    $st1 = "md5="
    $st2 = "denyip="
    $st3 = "filename="
    $st4 = "rmfile="
    $st5 = "exec_packet"
    $st6 = "build_iphdr"

  condition:
    all of them
}
```

Figure 6: Yara Rule

Here we used strings noticed in the binary itself:

```
string $st0 - ASCII XOR key
strings $st1 - $st4 are parameters used by functionalities in the binary related to trojan activities.
string $st5 - function name to execute DDoS flood
string $st6 - function name to build IP header
```

To match SYN Flood attack traffic generated by this botnet you can use the following tcpdump filters:

```
'ip[2:2] > 0x50' and 'tcp[13] & 2 !=0' and 'tcp[14:2] == 0xffff' and 'tcp[20:4] == 0x020405b4' and 'tcp[24:4] == 0x01010402'
```

Figure 7: Example tcpdump filters for SYN flood attack traffic

This filter is composed of the following BPF filters:

```
'ip[2:2] > 0x50' - matches on total size field greater than 80 bytes
'tcp[13] & 2 !=0' - matches on SYN flags
'tcp[14:2] == 0xffff' - matches on TCP Window size 65535
'tcp[20:4] == 0x020405b4' - matches on static bytes used for TCP header options
'tcp[24:4] == 0x01010402' - matches on static bytes used for TCP header options
```

To match DNS Flood attack traffic generated by this botnet you can use the following tcpdump filters:

```
'udp[10:4] == 0x01200001' and 'udp[14:4] == 0x00000000' and 'udp[19] == 0x0001' and 'ip[((ip[2:2]) - 14):1] = 0x01'
```

Figure 8: Example tcpdump filters for DNS flood attack traffic

The above BPF filter is composed of:

```
'udp[10:4] == 0x01200001' - matches on the specific DNS fields from the outlisted characteristics in Section 3
'udp[14:4] == 0x00000000' - matches on Answer Count and NS (Auth) Count fields
'udp[19] == 0x0001' - matches on Additional Record Count
'ip[((ip[2:2]) - 14):1] = 0x01' - matches on the DNS Query Type field for 'A' records
```

To remove the malware:

- Identify the main PID and the dropped files in /boot and /etc/init.d

```
root@ubuntu:/# ls -lha /boot | egrep " [a-z]{10}$"
-rwxr-x--- 1 root root 606K Aug 19 11:42 utmhsahoca

root@ubuntu:/# ls -lha /etc/init.d/utmhsahoca
-rwxr-x--- 1 root root 28 Aug 19 11:42 /etc/init.d/utmhsahoca

root@ubuntu:/# find /proc/ -name exe 2> /dev/null | xargs -I{} ls -l {} 2> /dev/null | egrep
'/boot/[a-z]{10}$'
lrwxrwxrwx 1 root root 0 Aug 19 11:42 /proc/3746/exe -> /boot/utmhsahoca
```

Figure 9: Bash script used in the /proc file system process

The above bash script looks through all of the processes in the /proc file system and gets the path of their symlinked source finding all processes started from an executable in /boot with a 10-character name.

Here we have identified that /boot/utmhsahoca is currently running user PID 3746.

The above script in more readable form and commented form:

```
# find all executing processes
1. find /proc/ -name exe 2> /dev/null

# get each processes symlinked destination
2. xargs -I{} ls -l {} 2> /dev/null

# find symlinks that match our '/boot/xxxxxxxxx'
signature
3. egrep '/boot/[a-z]{10}$'
```

Figure 10: Additional Bash script

```
root@ubuntu:/boot# ps -eaf -u root | grep `date +%H:%M` | egrep -
v "ps|grep"
root      8318   2868   0 12:04 ?        00:00:00 ps -
ef
root      8321   2868   0 12:04 ?        00:00:00 cat
resolv.conf
root      8324   2868   0 12:04 ?        00:00:00 ls -
la
root      8326   2868   0 12:04 ?        00:00:00
id
root      8327   2868   0 12:04 ?        00:00:00 ls -
la
```

Figure 11: Identification for the supporting processes responsible for the persistence of the main process

The PIDs 8318, 8321, 8324, 8326, 8327 are short lived and their purpose is to only check if the main process is running. We are going to eliminate them by their start time, which will always be very close to the current time. If this command is run multiple times over the span of a minute the observed PIDs and associated fake commands will constantly change.

### 3.1 / KILLING THE MALICIOUS PROCESS /

```
root@ubuntu:/boot# ps -eaf | grep ^root | grep `date +%H:%M` | egrep -v "ps|grep" | awk '{print
$2}' | xargs -I{} kill -9 {}; kill -9 3746
```

Figure 12: Bash script in more readable form

```

# get processes
1. ps -eaf

# only get root users processes
2. grep ^root

# only keep processes started within the current minute
3. grep `date +%H:%M`

# remove grep and ps from list from the list
4. egrep -v "ps|grep"

# get the PID of the process
5. awk '{print $2}'

# force-kill the PID supplied for subprocesses
6. xargs -I{} kill -9 {};

# kill the main process
7. kill -9 3746

```

Figure 13: Additional Bash script

This script will get all processes running under the root user, it will then grab a subset of those processes that have spawned within the last minute, removing our own `ps` and `grep` from the list. It then extracts the PID for each process and kills it using `-9`, which cannot be blocked, prevented, or ignored. Once it has killed all of the watcher processes it kills off the master process, in this case PID **3746**.

Next we'll need to delete all of the malware files that have been dropped to disk.

```

root@ubuntu:/boot# rm -f /boot/utmhsahoca && rm -f /etc/init.d/utmhsahoca && /lib/udev/udev

```

Figure 14: Delete the dropped files in `/boot` and `/etc/init.d`

If you failed to kill the running instances of the malware, after step 4 the malware will drop a new copy of itself with a new 10-character randomized name in `/boot` and new startup script in `/etc/init.d`. If new files exist repeat steps 1-4 to ensure you've killed all running instances. You can also use the above YARA rule (Figure 6) on the whole file system to see if any copies of the malware are still present.



**4.0 / CONCLUSION** / Akamai's SIRT expects XOR DDoS activity to continue as attackers refine and perfect their methods.

This will likely result in a more diverse selection of DDoS attack types included in future versions of the malware. XOR DDoS malware is part of a wider trend of which companies must be aware: Attackers are targeting poorly configured and unmaintained Linux systems for use in botnets and DDoS campaigns.

A decade ago, Linux was seen as the more secure alternative to Windows environments, which suffered the lion's share of attacks at the time, and companies increasingly adopted Linux as part of their security-hardening efforts.

As the number of Linux environments has grown, the potential opportunity and rewards for criminals has also grown. Attackers will continue to evolve their tactics and tools and security professionals should continue to harden their Linux based systems accordingly.



**About Akamai Security Intelligence Response Team (SIRT)** Focuses on mitigating malicious global cyber threats and vulnerabilities, the Akamai Security Intelligence Response Team (SIRT) conducts and shares digital forensics and post-event analysis with the security community to proactively protect against threats and attacks. As part of its mission, the Akamai SIRT maintains close contact with peer organizations around the world and trains Akamai's Professional Services and Customer Care teams to both recognize and counter attacks from a wide range of adversaries. The research performed by the Akamai SIRT is intended to help ensure Akamai's cloud security products are best of breed and can protect against any of the latest threats impacting the industry.

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